



Forum on specification & Design Languages (FDL'20)



FDL stimulates scientific and controversial discussions in a friendly and productive environment.

New trends and traditional topics in the broad fields of embedded/electronics/software systems and languages merge in a lively and cross-discipline research & industrial community.

Calls for Special sessions, Full (8 pp), short (4 pp), and WiP/PhD Forum/Poster (2 pp) papers.



Keynotes: Edward Lee / UC Berkeley,
Manuel Serrano / Inria & Université Côte d'Azur,
Hauke Fuhrmann / Scheidt & Bachmann



7–9 September 2020 | Kiel, Germany



Deadlines:

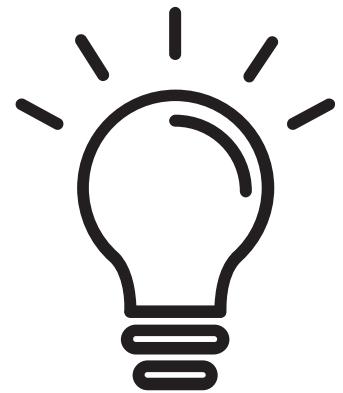
Special Sessions:	March 22, 2020
Paper Deadline:	May 29, 2020
PhD/WiP Deadline:	June 12, 2020
Author Notification:	June 28, 2020
Final Version:	July 19, 2020

Website: www.fdl-conference.org | Contact: fdl2020@easychair.org

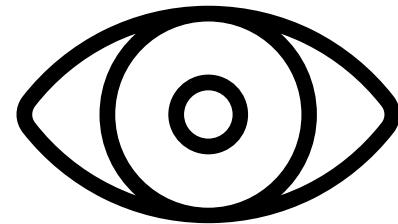
Re FDL'19: Open call for ACM TECS Special Issue on Specification and Design Languages
Deadline: Feb. 1, 2020 (firm)
Contacts: Alain Girault, Reinhard von Hanxleden

From Lustre to Graphical Dataflow Models

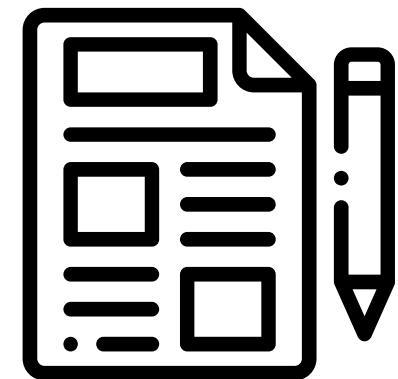
Lena Grimm
Reinhard von Hanxleden
Kiel University



Motivation



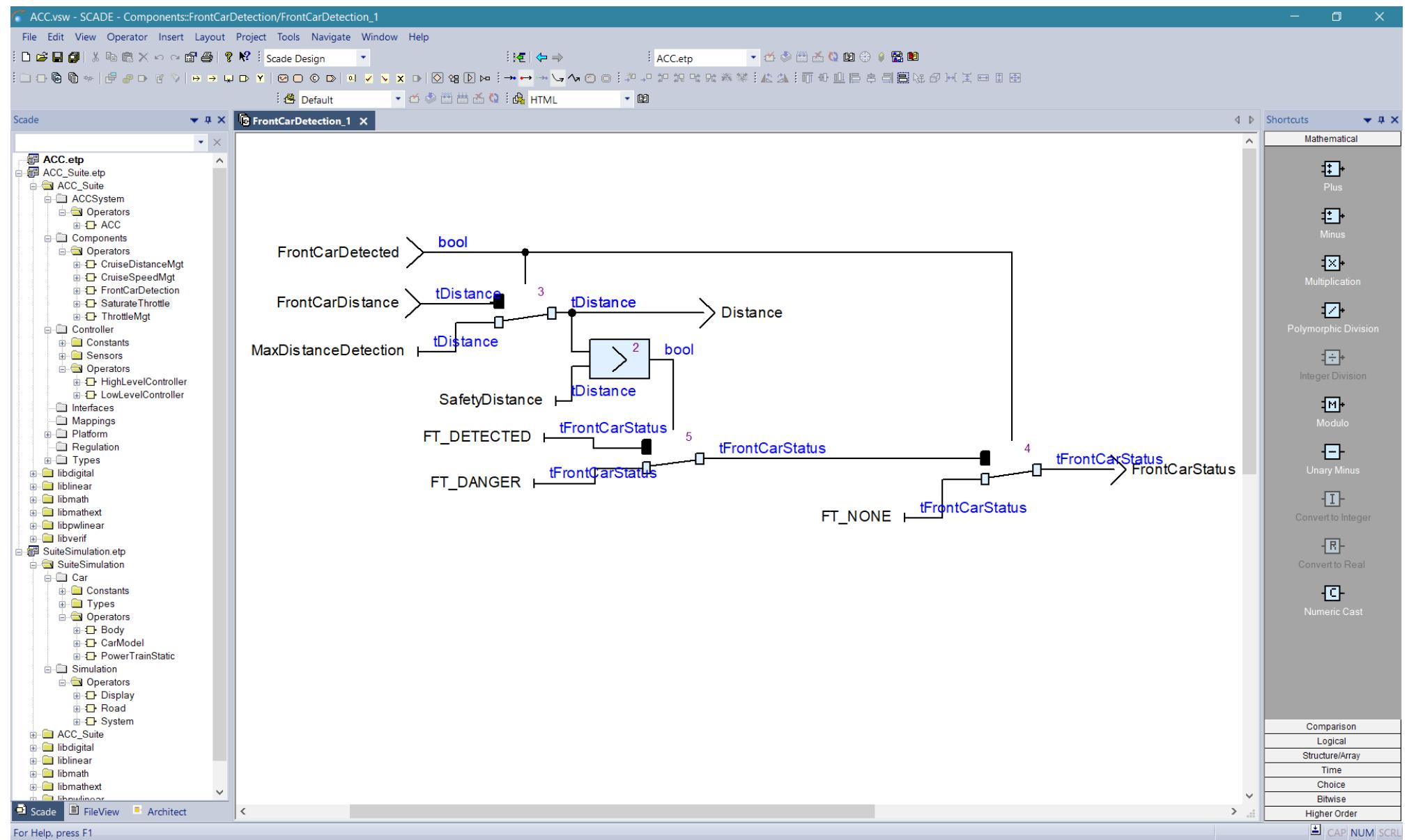
Visual



Concepts



SCADE





KIELER

File Edit Navigate Search Project Run Window Help

dataflow.sctx

```
1 scchart dataflowChart {
2
3     input bool in1, in2
4     output bool o
5
6     dataflow {
7         o = !in1 & in2
8     }
9 }
```

Editing

Diagram

dataflowChart

input bool in1, in2
output bool o

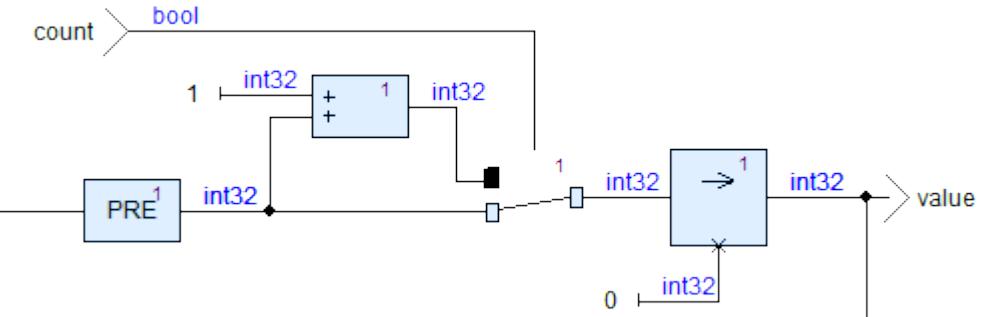
Automatic Diagram Generation

Writable Insert 9:2:97



SCADE

```
L1 = count;
value = L7;
L2 = 1;
L3 = pre L7;
L4 = L2 + L3;
L5 = if L1 then (L4) else (L3);
L6 = 0;
L7 = (L6) -> (L5);
```

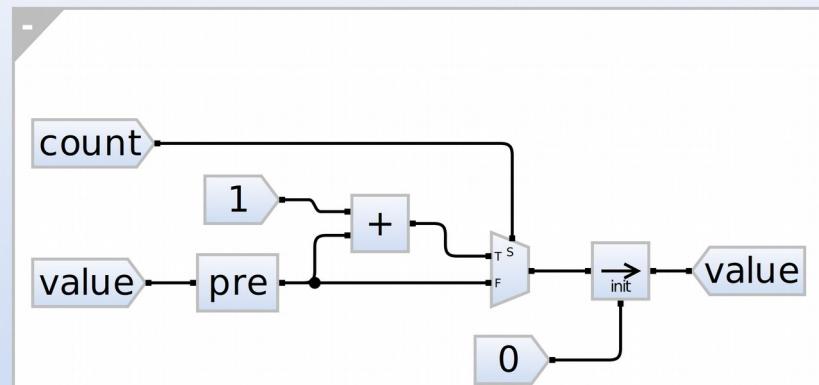


```
node Counter (count:bool) returns (value:int)
var L1:bool;
L2, L3, L4, L5, L6, L7:int;
let
  L1 = count;
  value = L7;
  L2 = 1;
  L3 = pre L7;
  L4 = L2 + L3;
  L5 = if L1 then (L4) else (L3);
  L6 = 0;
  L7 = (L6) -> (L5);
tel
```

KIELER

Counter

input bool count
output int value





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Editor Support for Lustre

Automatic Diagram Recovery

Counter

input bool count
output int value

The diagram illustrates the internal logic of the Counter node. It takes a boolean input 'count' and an integer input 'value'. If 'count' is true, the 'value' is incremented by 1 using a sum node ('+') and the result is stored in a variable 'pre'. If 'count' is false, the variable 'pre' remains unchanged. The initial value of 'pre' is 0, which is set by an 'init' block. The final value is then output as 'value'.

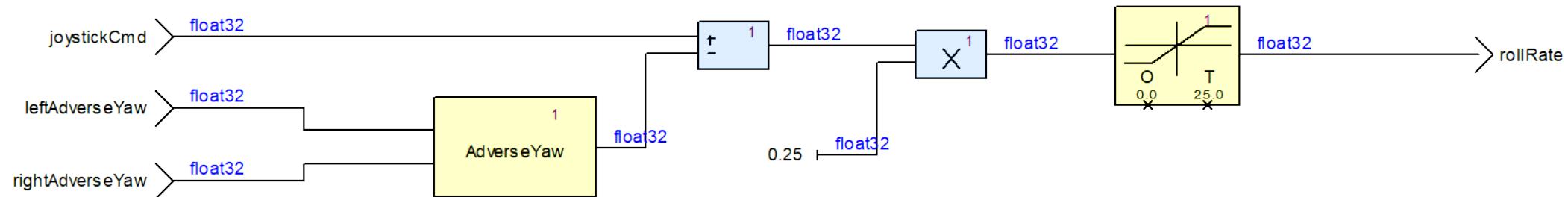
```
1 node Counter (count : bool)
2   returns (value : int)
3
4 var preValue:int;
5 let
6   preValue = pre value;
7   value = 0 ->
8     if count
9       then 1 + preValue
10      else preValue;
11 tel
```



Demo

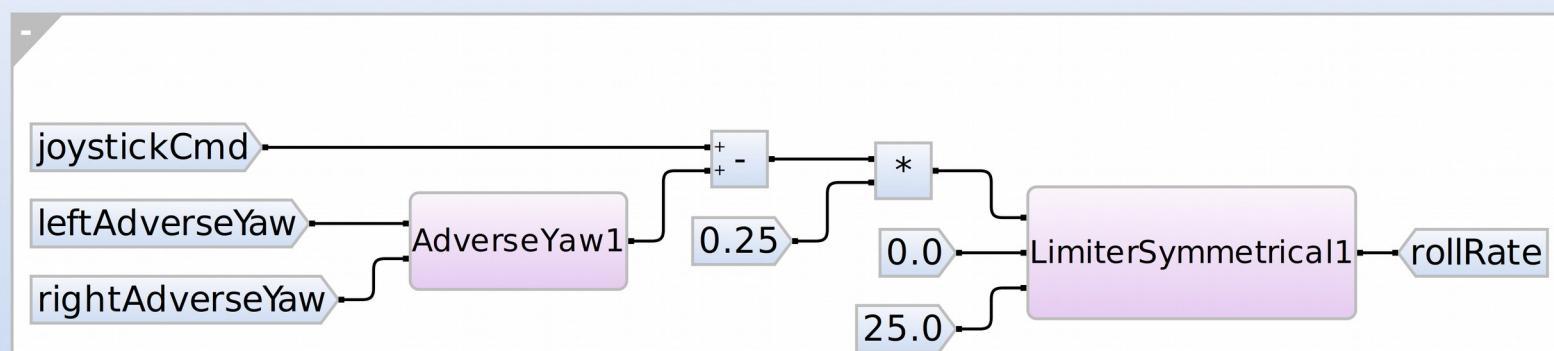


Larger Example



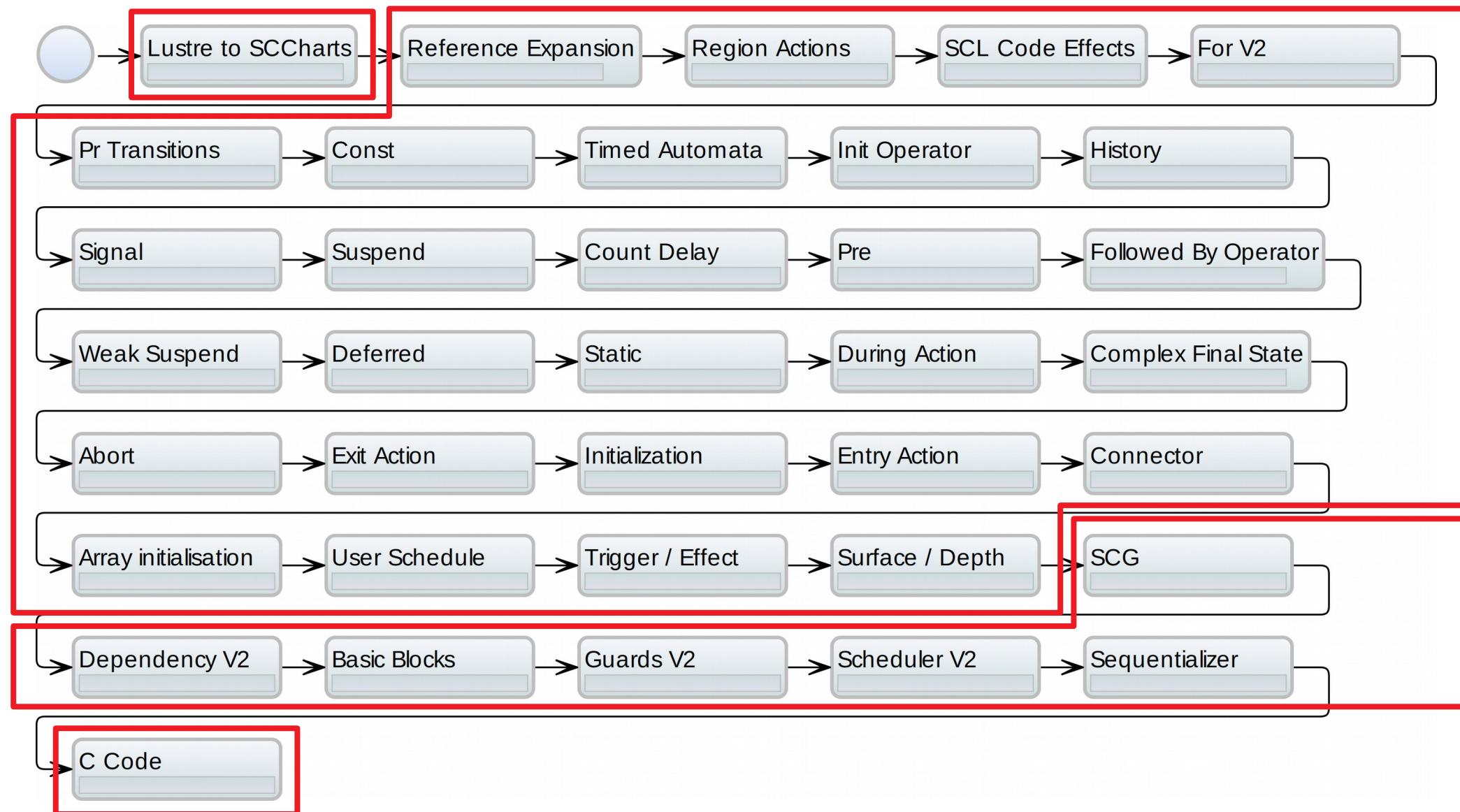
RollRateCalculate

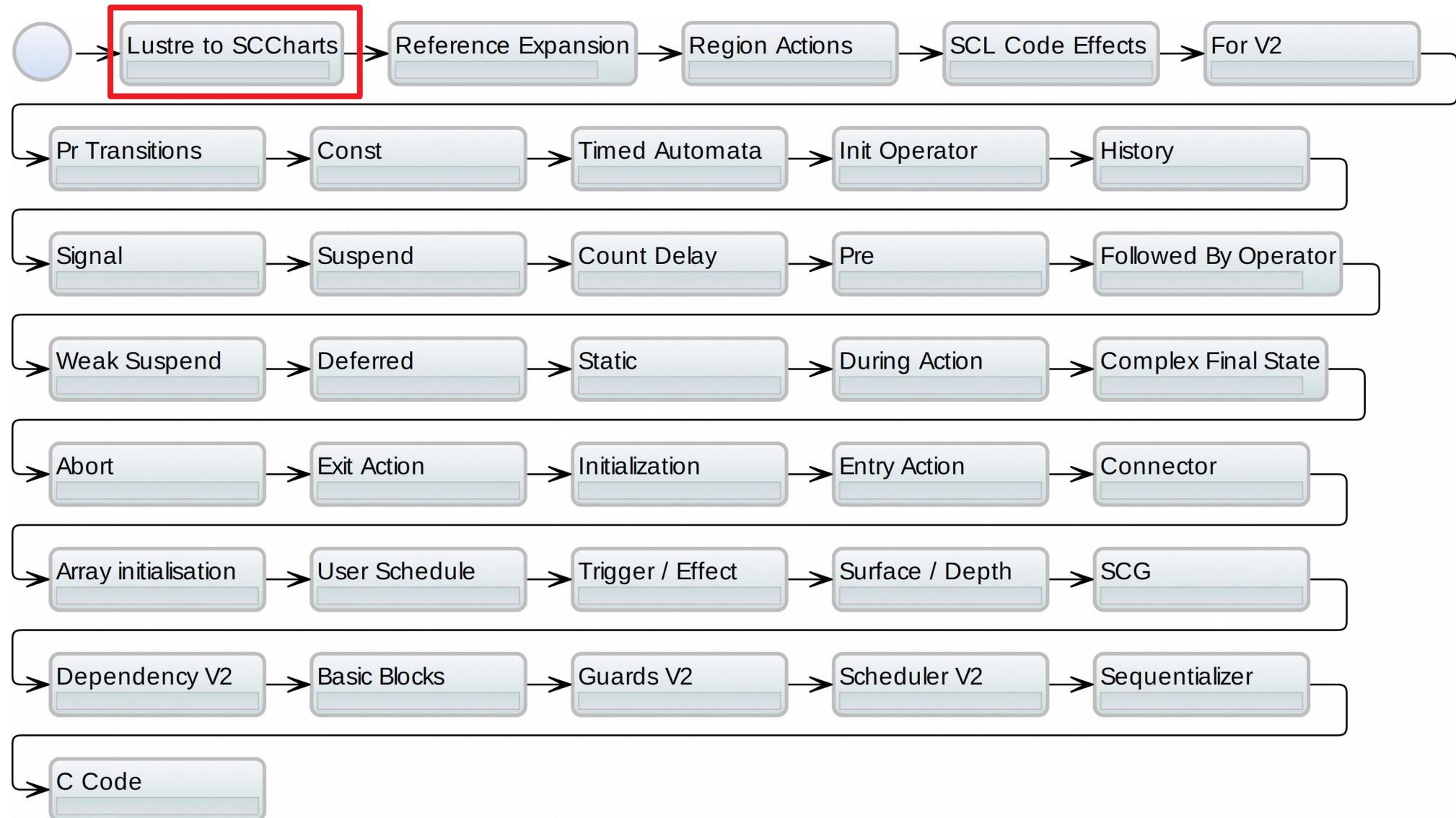
```
input float joystickCmd
input float leftAdverseYaw
input float rightAdverseYaw
output float rollRate
ref AdverseYaw AdverseYaw1
ref LimiterSymmetrical LimiterSymmetrical1
```





KIELER Compilation Chain







Transformation Challenges

Lustre

x	1	2	3	4	5	6	7	8	9
clk	true	false	true	false	false	true	false	true	true
x when clk	1		3			6		8	9



when, current?

SCCharts

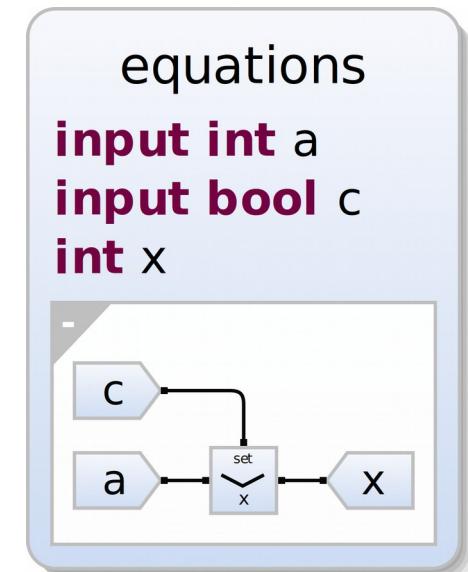
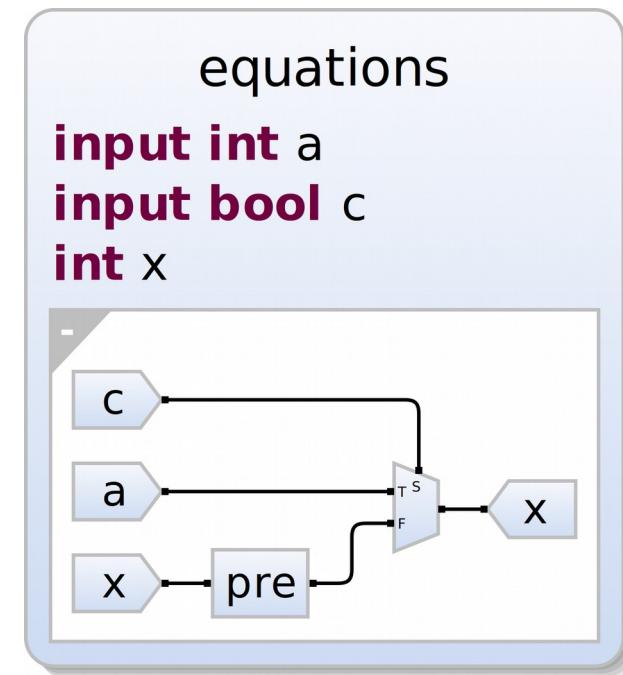
x	1	2	3	4	5	6	7	8	9
clk	true	false	true	false	false	true	false	true	true
clk? x	1	1	3	3	3	6	6	8	9



when Operator

```
node equations(a:int; c:bool)
  returns ();

var x:int when c;
let
  x = a when c;
tel.
```



```
scchart equations {
  input int a
  input bool c
  int x

  dataflow {
    x = c ? a : pre(x)
  }
}
```

```
scchart equations {
  input int a
  input bool c
  int x

  dataflow {
    x = c ? a
  }
}
```



when Operator

clk	true	false	true	false	true	true	false	false	true
x	true	false	false	true	true	false	false	false	true
y	true	false	false	true	false	false	true	true	false

xClk = x when clk	true	false	true	false	true
y when xClk	true		false		false



When Operator with Variables I

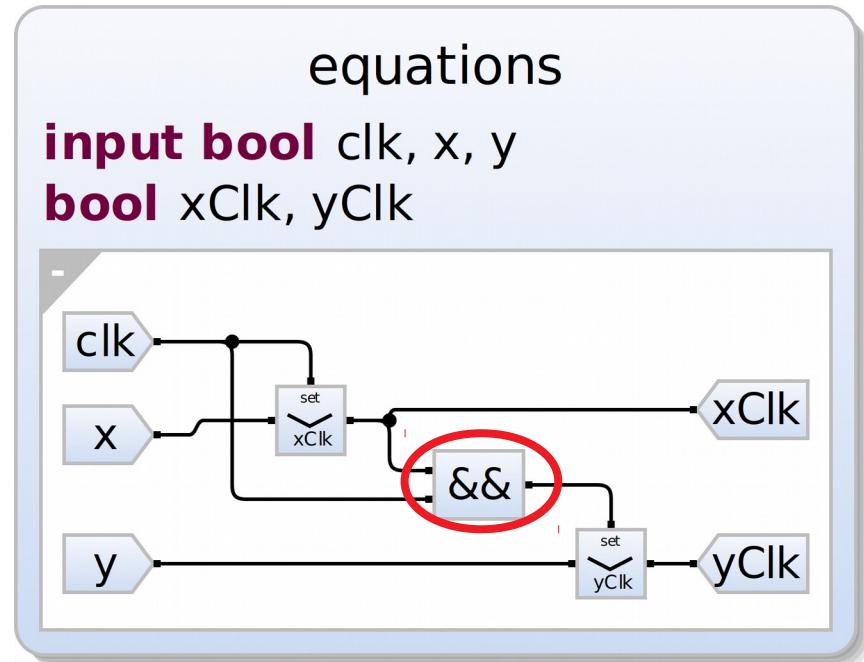
clk	true	false	true	false	true	true	false	false	true
x	true	false	false	true	true	false	false	false	true
y	true	false	false	true	false	false	true	true	false

xClk = clk? x	true	true	false	false	true	false	false	false	true
xClk? y	true	false	false						



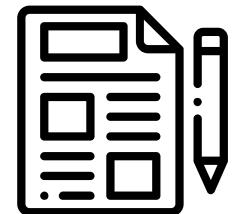
when Operator

```
node equations(clk,x,y:bool)
  returns ();
var xClk:bool when clk;
  yClk:bool when xClk;
let
  xClk = x when clk;
  yClk = y when xClk;
tel.
```



```
scchart equations {
  input bool clk, x, y
  bool xClk, yClk

  dataflow {
    xClk = clk ? x
    yClk = (xClk && clk) ? y
  }
}
```



When Operator with Variables II

clk	true	false	true	false	true	true	false	false	true
x	true	false	false	true	true	false	false	false	true
y	true	false	false	true	false	false	true	true	false

xClk = clk? x	true	true	false	false	true	false	false	false	true
(clk && xClk)? y	true	true	true	true	false	false	false	false	false





Current

clk	true	false	true	false	true	true	false	false	true
x	true	false	false	true	true	false	false	false	true
y	true	false	false	true	false	false	true	true	false

xClk = clk? x	true	true	false	false	true	false	false	false	true
(clk && xClk)? y	true	true	true	true	false	false	false	false	false

Always implicit ‘current’ through Variables



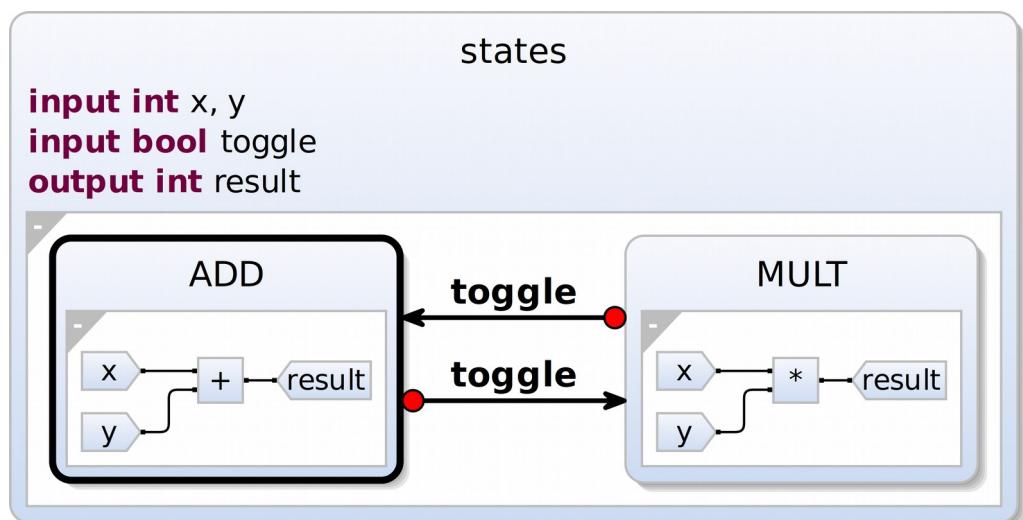


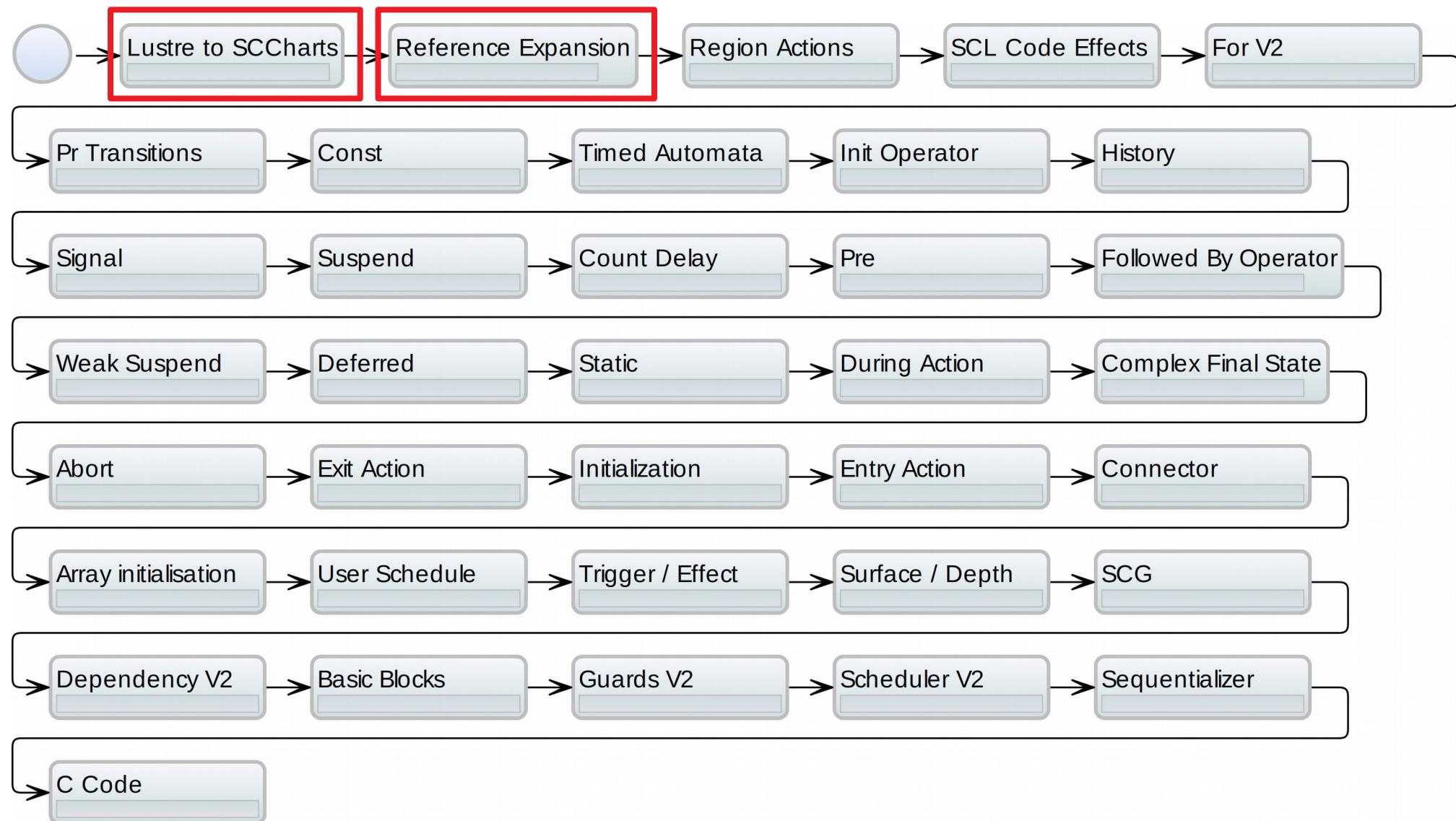
State Extension

```
node states (x,y:int; toggle:bool)
    returns (result:int);
let
    automaton Calculator
        state ADD
        let
            result = x + y;
        tel
        unless if toggle restart MULT;

        state MULT
        let
            result = x * y;
        tel
        unless if toggle restart ADD;

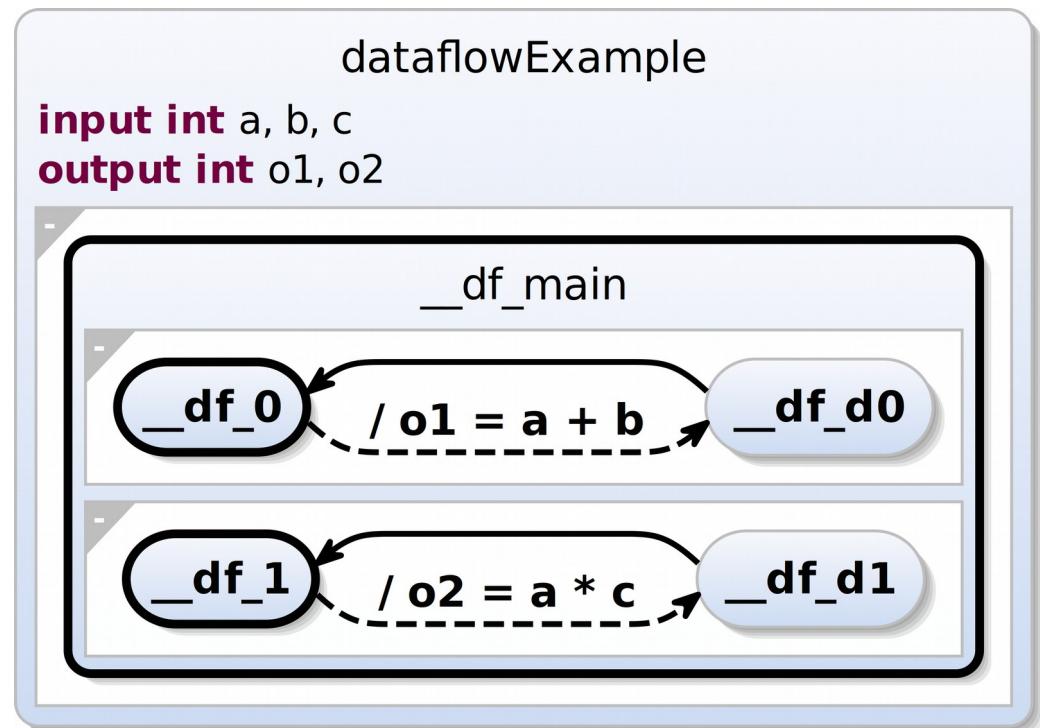
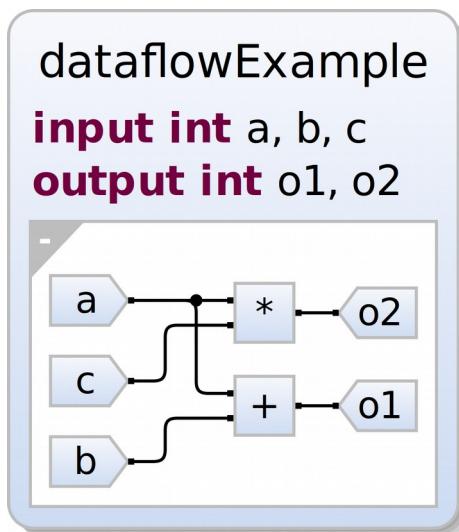
    returns ...
tel
```

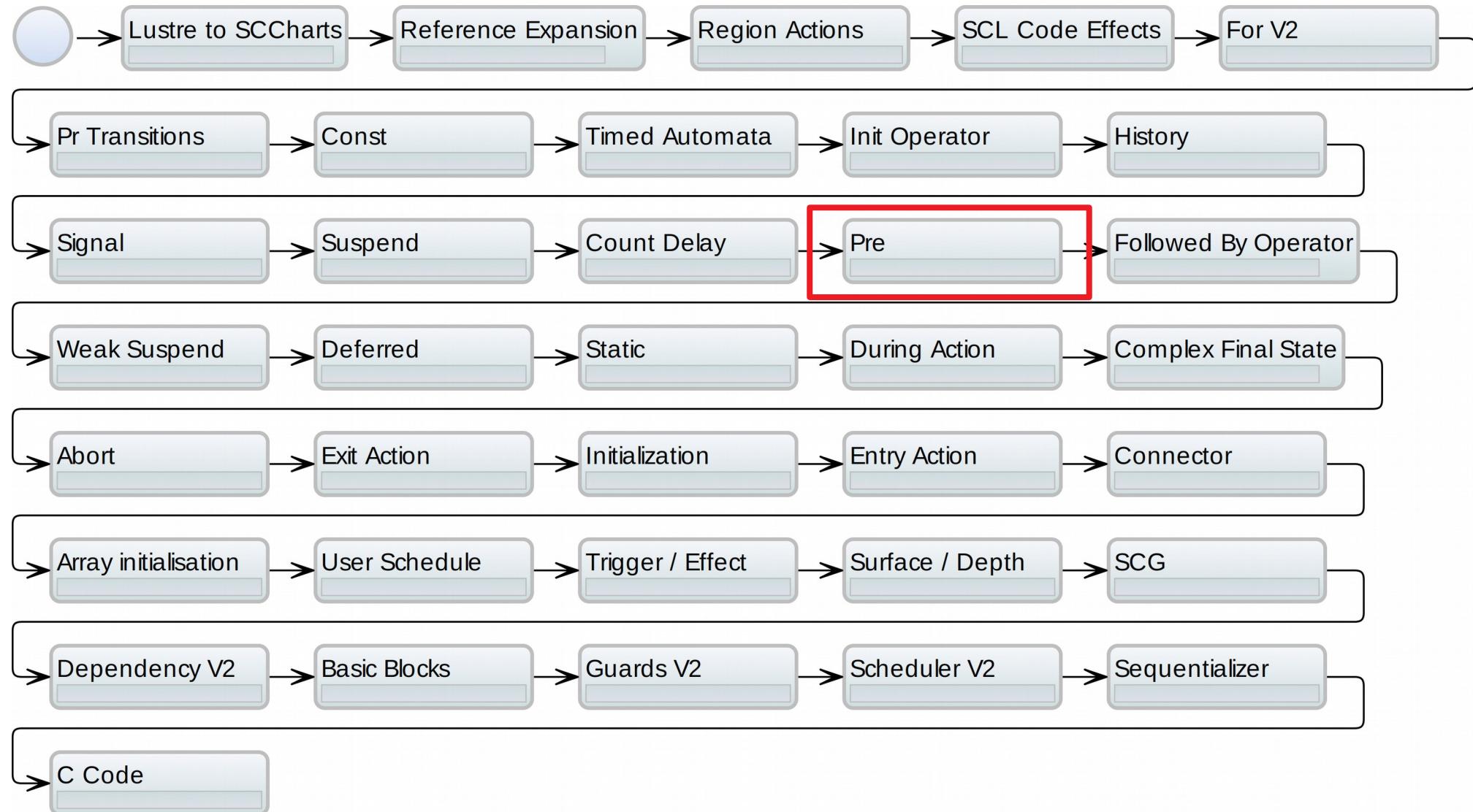






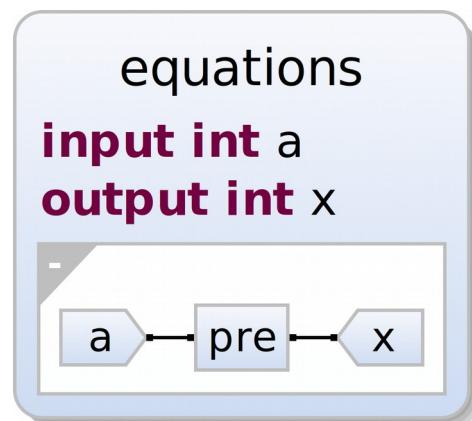
SCCharts Dataflow Semantics





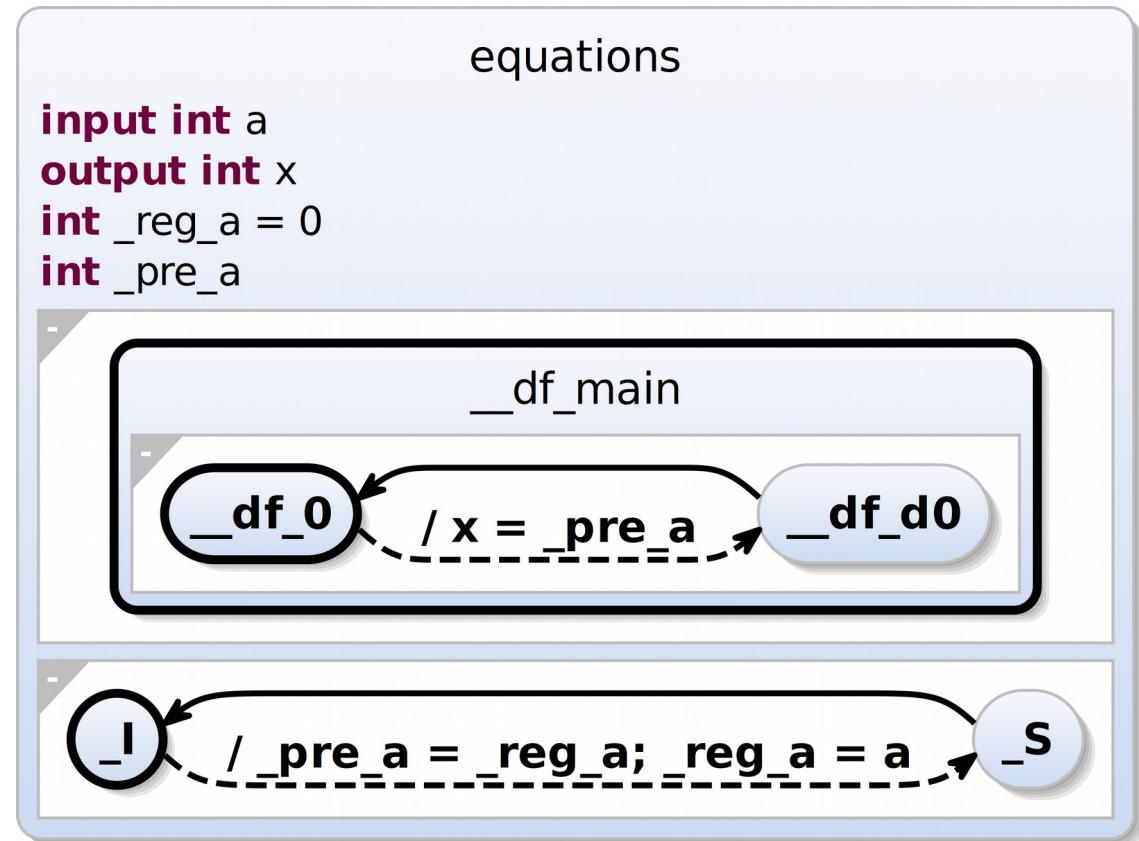


Pre Operator in SCCharts



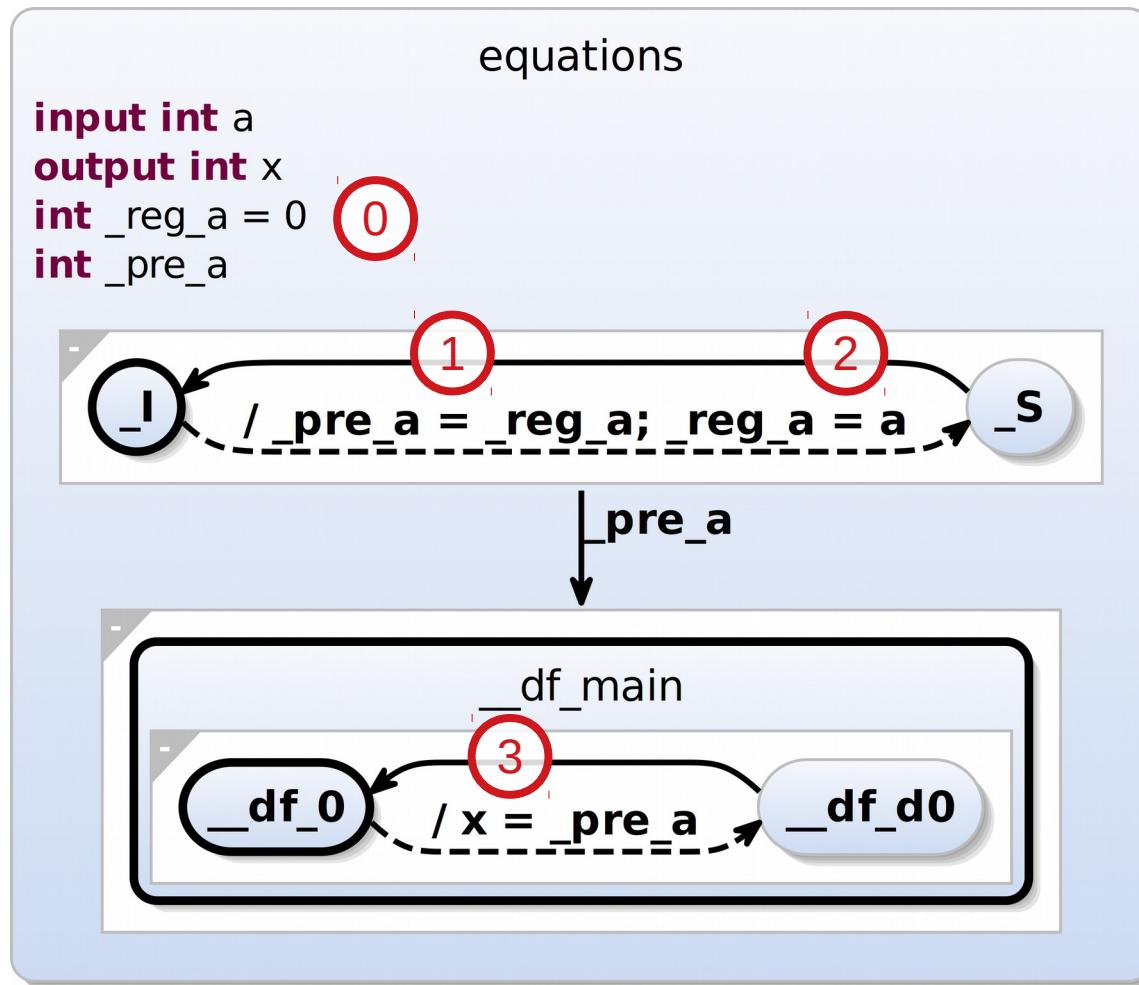
```
scchart equations {
    input int a
    output int x

    dataflow {
        x = pre(a)
    }
}
```





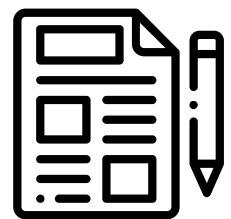
Pre Operator in SCCharts Induced Dataflow View





Pre Operator with Clocks

clk	true	false	false	true	true	false	true	false	true
x	1	2	3	4	5	6	7	8	9
xClk = x when clk	1			4	5		7		9
pxClk = pre(xClk)	nil			1	4		5		7
pre(pxClk)	nil			nil	1		4		5



Pre Operator with Clocks and Variables

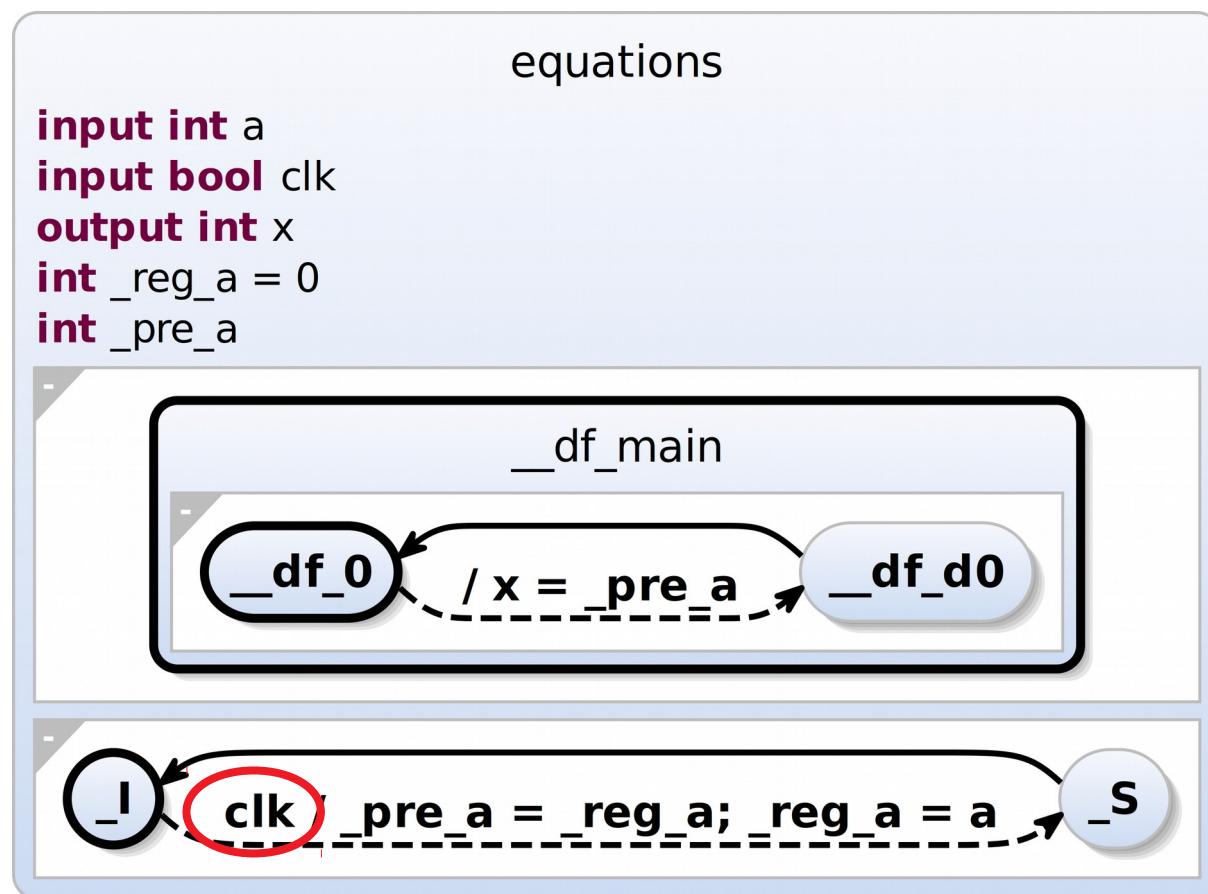
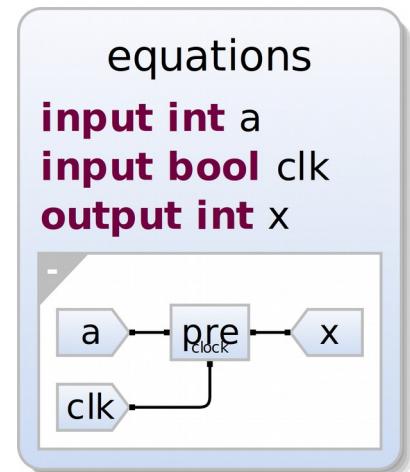
clk	true	false	false	true	true	false	true	false	true
x	1	2	3	4	5	6	7	8	9
xClk = clk? x	1	1	1	4	5	5	7	7	9
pxClk = pre(xClk)	nil	1	1	1	4	5	5	7	7
pre(pxClk)	nil	nil	1	1 nil	1	4	5 4	5	7 5

Diagram illustrating the state transitions and the application of the Pre operator:

- The first row shows the initial state of variables: clk, true, false, false, true, true, false, true, false, true.
- The second row shows the value of variable x at each step: 1, 2, 3, 4, 5, 6, 7, 8, 9.
- The third row shows the value of xClk = clk? x at each step: 1, 1, 1, 4, 5, 5, 7, 7, 9.
- The fourth row shows the value of pxClk = pre(xClk) at each step: nil, 1, 1, 1, 4, 5, 5, 7, 7.
- The fifth row shows the value of pre(pxClk) at each step: nil, nil, 1, 1
nil, 1, 4, 5
4, 5, 7
5.
- Red arrows indicate the flow of values from the pxClk row to the pre(pxClk) row, showing how the Pre operator retrieves the previous value of pxClk.



Clocked Pre Operator in SCCharts



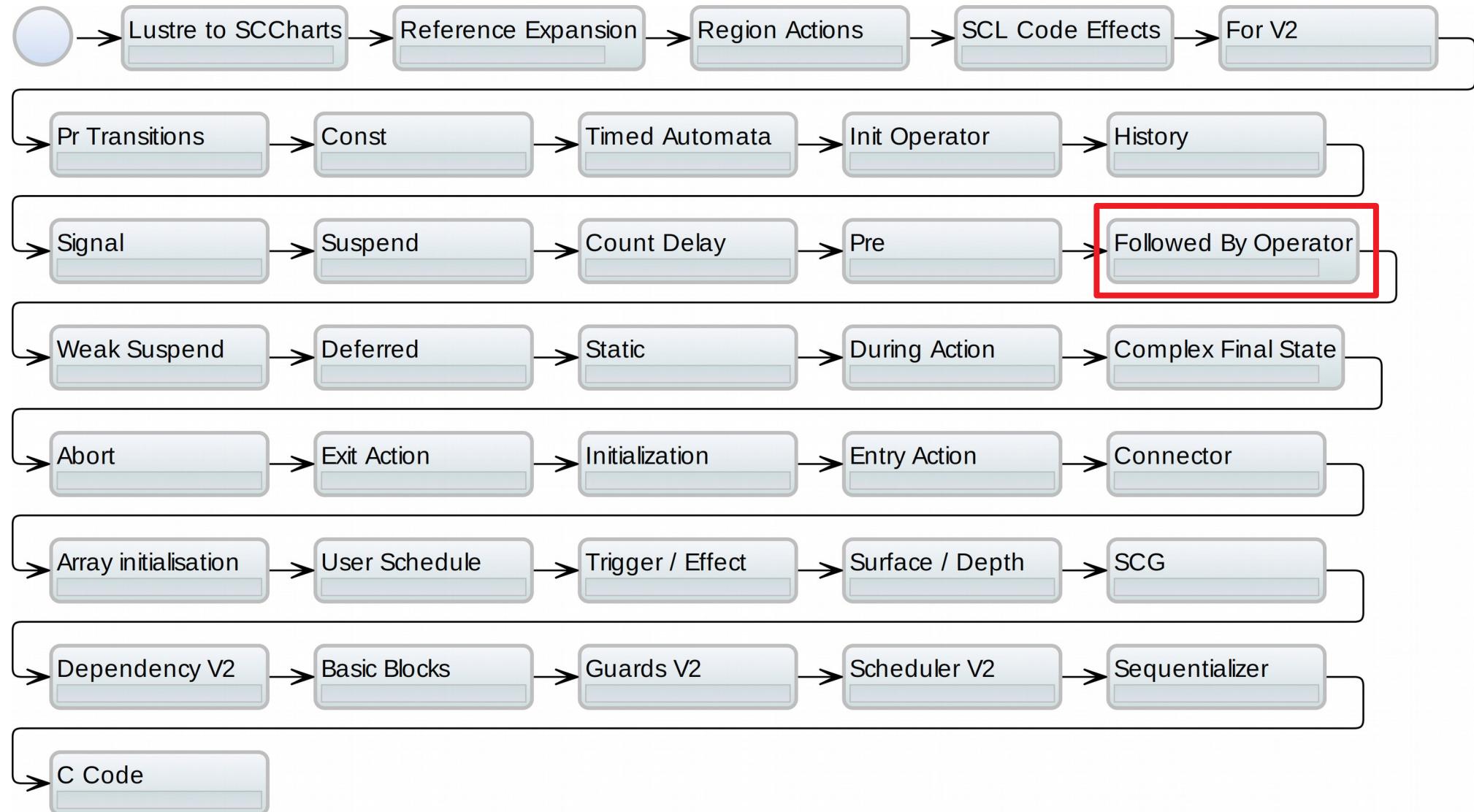
```
scchart equations {  
    input int a  
    input bool clk  
    output int x  
  
    dataflow {  
        x = pre(a, clk)  
    }  
}
```

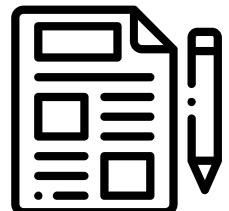


Pre Operator with Clocks and Variables

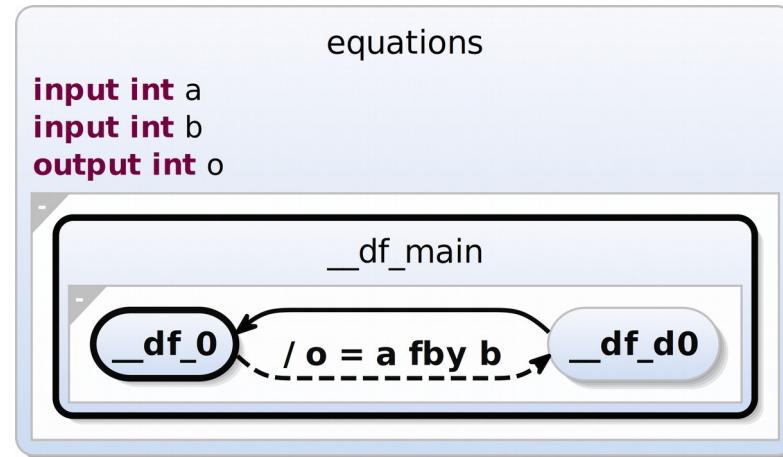
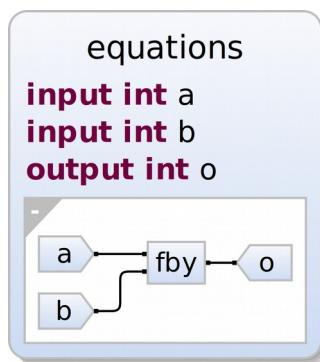
clk	true	false	false	true	true	false	true	false	true
x	1	2	3	4	5	6	7	8	9
xClk = clk? x	1	1	1	4	5	5	7	7	9
pxClk = pre(xClk, clk)	nil	nil	nil	1	4	4	5	5	7
pre(pxClk, clk)	nil	nil	nil	nil	1	1	4	4	5



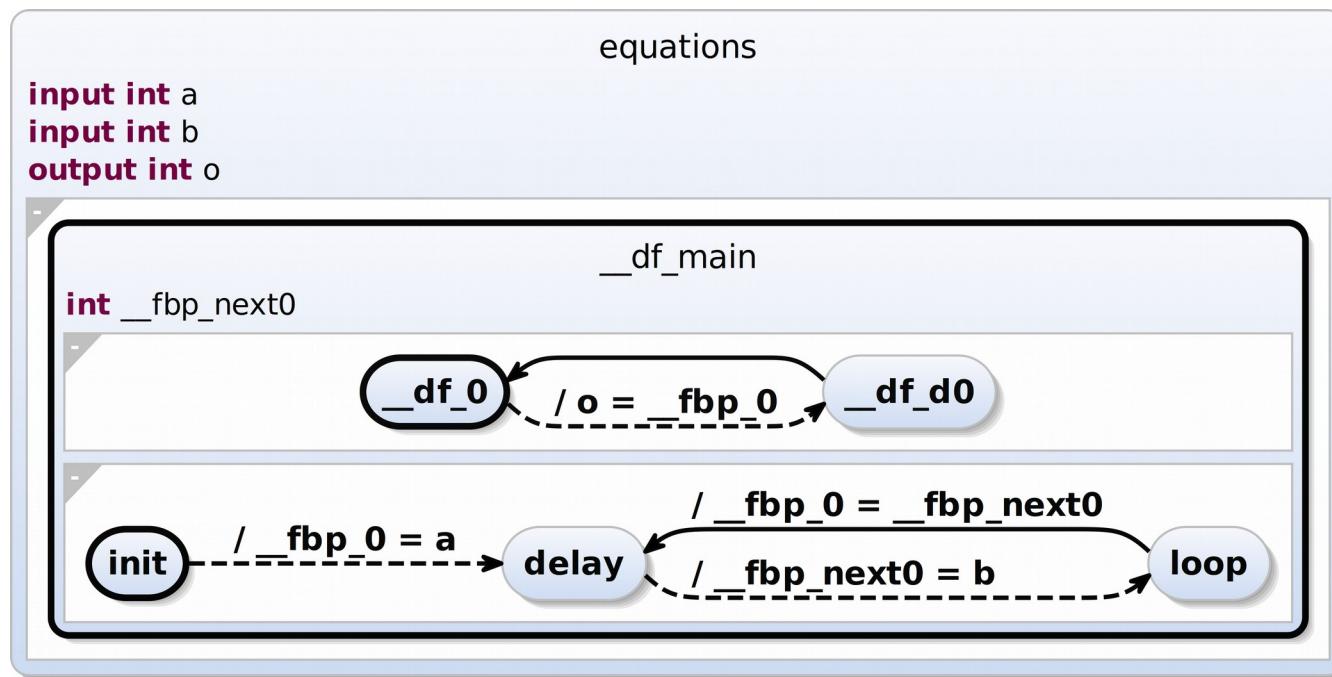




Fby Transformation

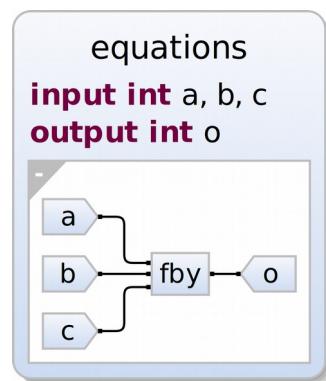


$a \rightarrow \text{pre}(b)$



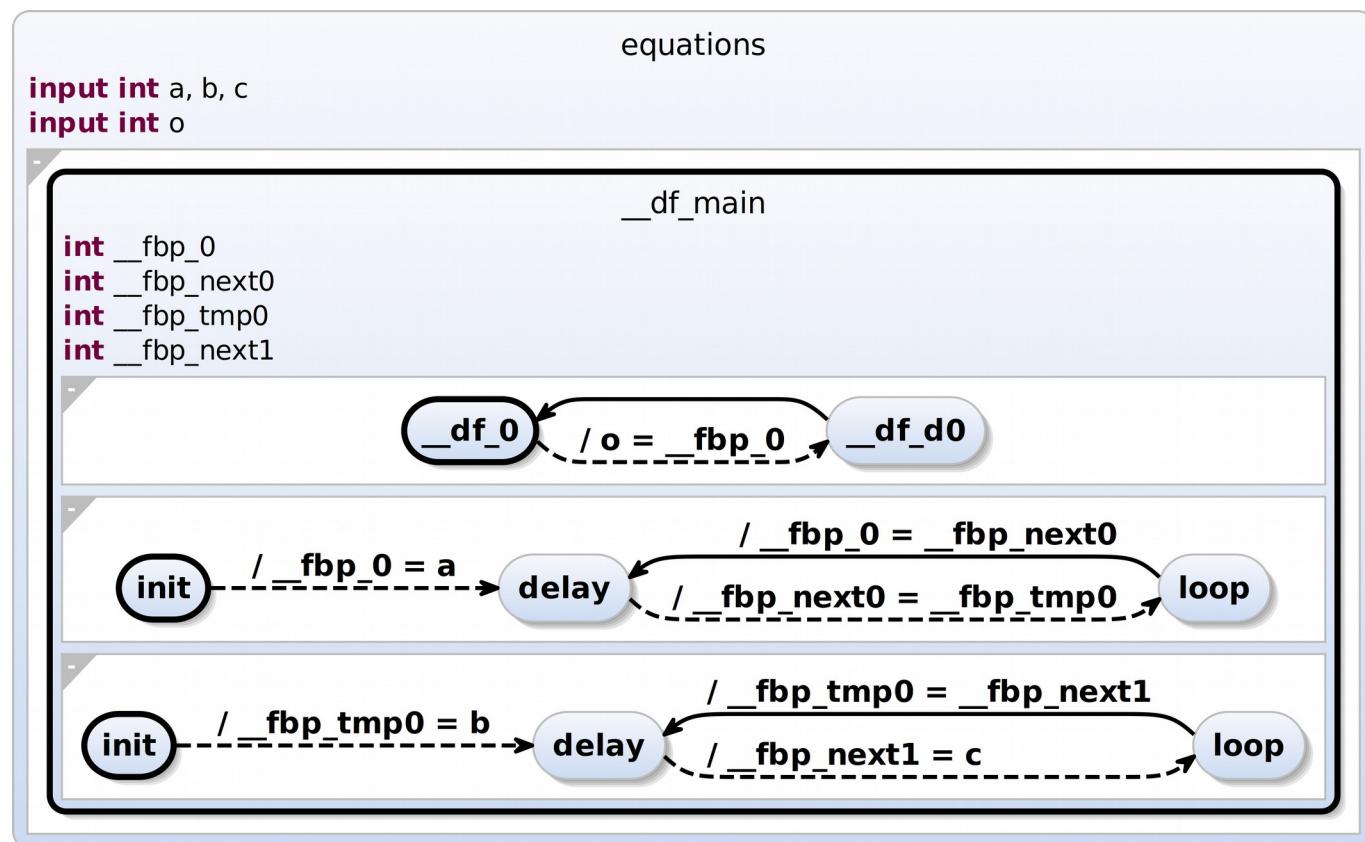


Fby Transformation II



```
scchart equations {
    input int a, b, c
    input int o

    dataflow {
        o = a fby b fby c
    }
}
```

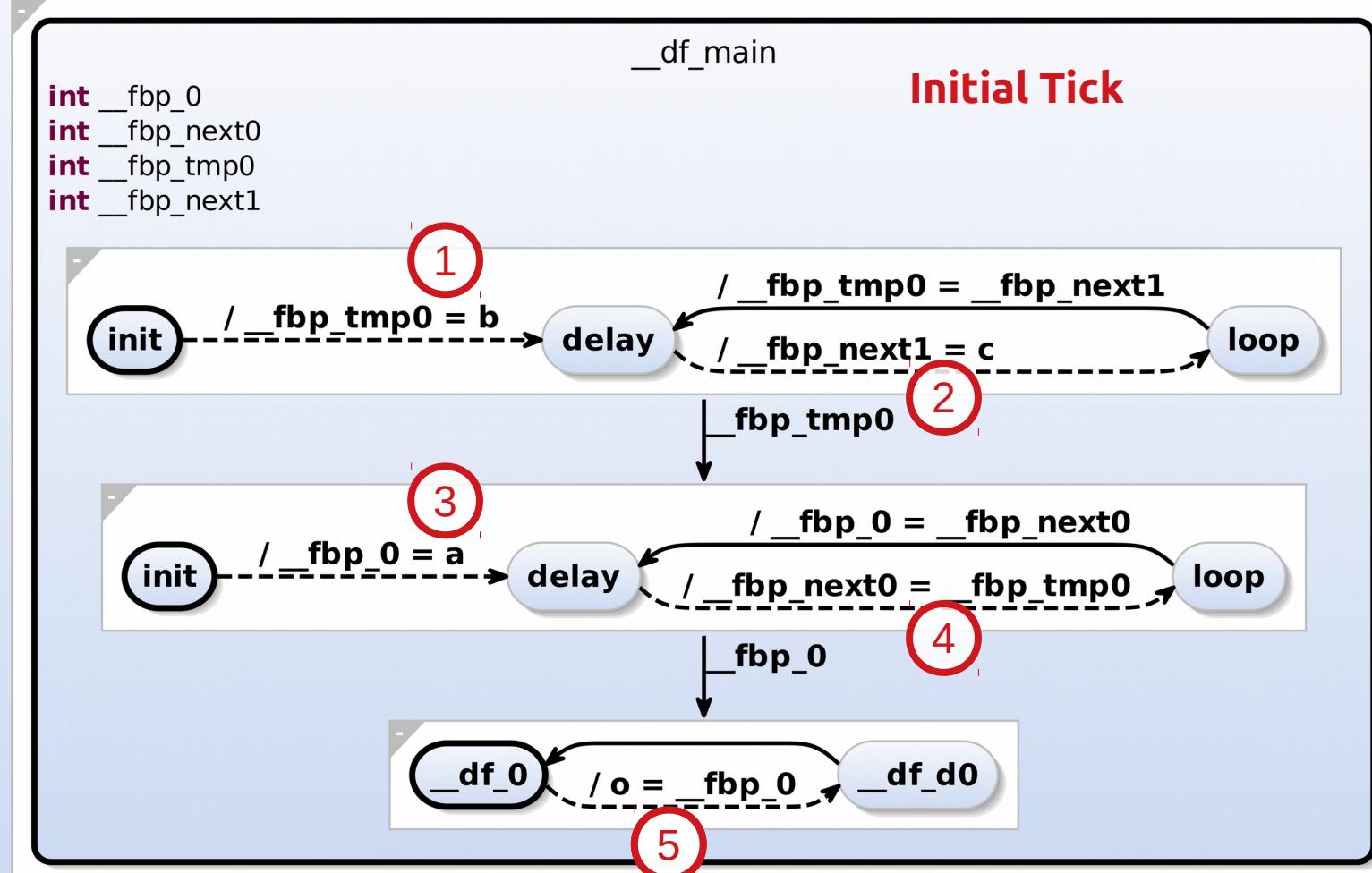




Fby Transformation Induced Dataflow View

equations

```
input int a, b, c  
input int o
```

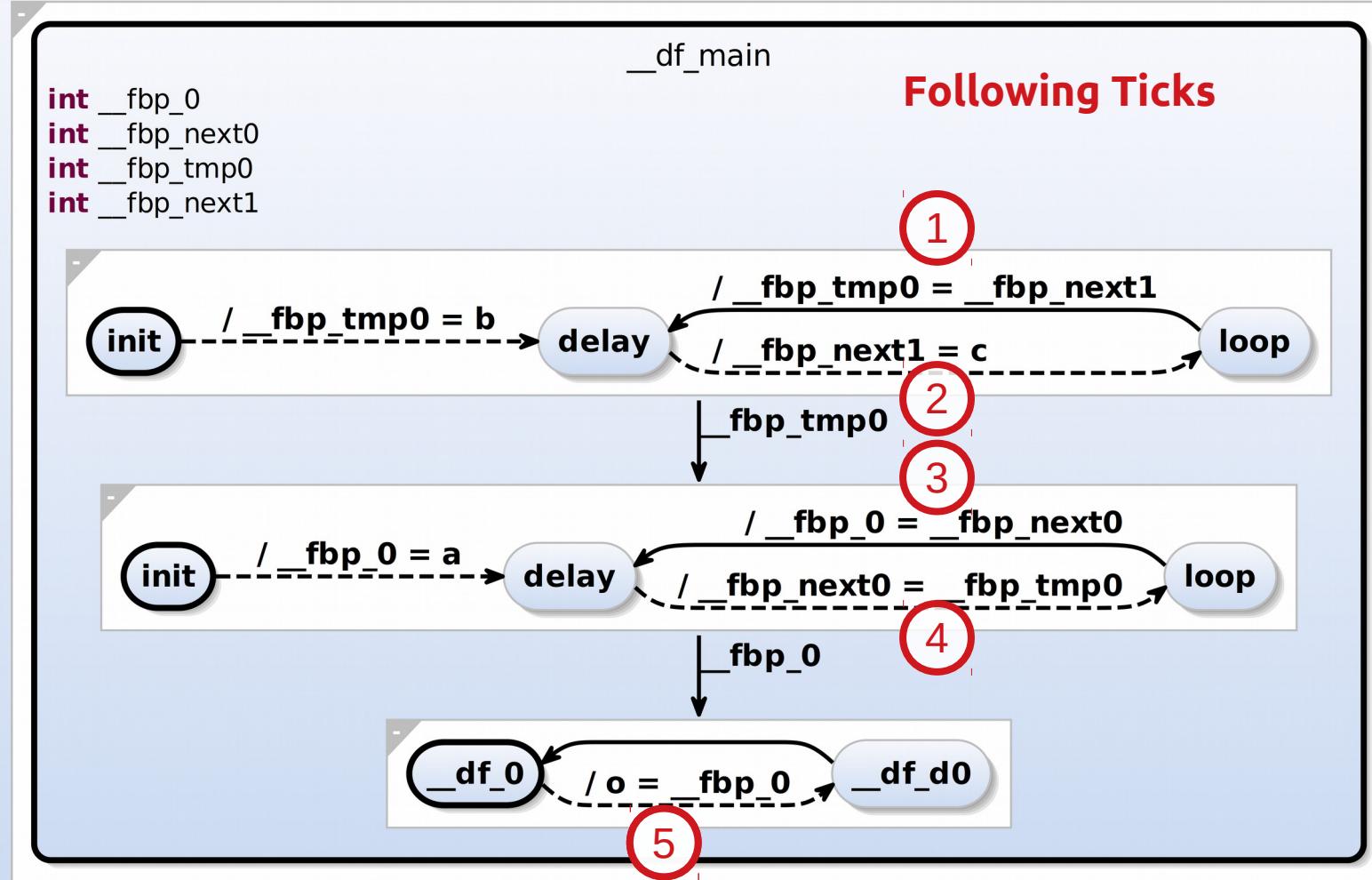




Fby Transformation Induced Dataflow View

equations

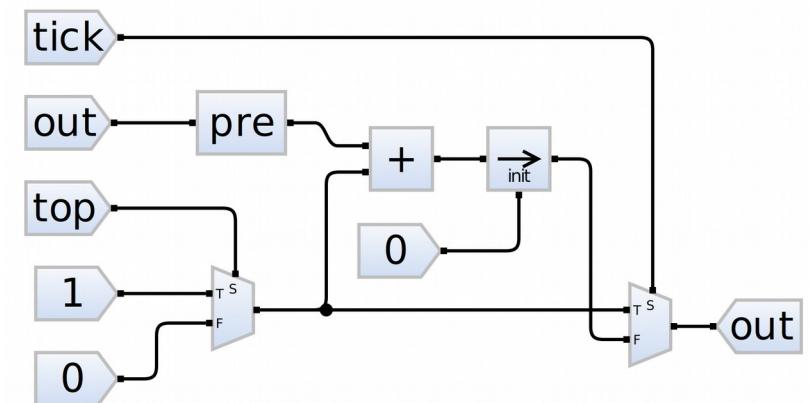
input int a, b, c
input int o



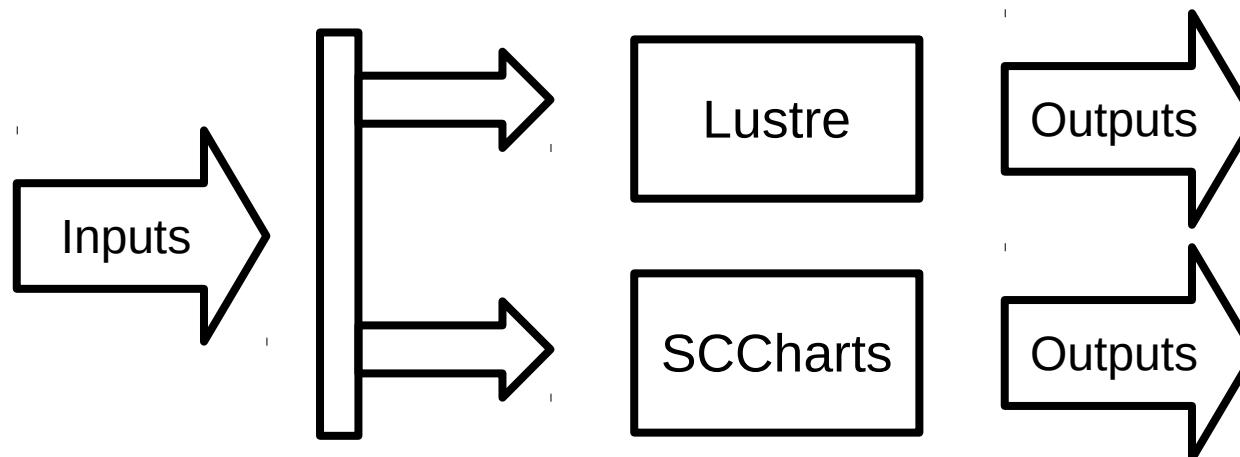
Summary

Model Recovery

```
node counting(tick:bool;top:bool)
    returns (out:int);
var v:int;
let
    v = if top then 1 else 0;
    out = if tick
        then v
        else (0 -> pre out + v);
tel.
```



Behavior Preservation



Thank You!

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<https://rtsys.informatik.uni-kiel.de/confluence/display/KIELER/Downloads>